# On the Estimation of Joint Mutual Information for Physical Layer Security

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Physical Layer Security



- ► The reciprocal channel (h<sub>a</sub> = h<sub>a'</sub>) can be used to generate the secret key.
- Requires channel to be changing with time (fading).

# Available Key bits $(I_{\rm K})$

- How many secret key bits can be generated per observation of the channel?
- Depends on the mutual information between the two channels.
- Mutual Information: Amount of information shared between  $\hat{h}_a$  and  $\hat{h}_{a'}$ . How much information  $\hat{h}_a$  tells us about  $\hat{h}_{a'}$  and vice versa.

$$I_{\rm K} = I(\hat{h}_a; \hat{h}_{a'}) = \mathbf{E} \left\{ \log_2 \frac{f(\hat{h}_a, \hat{h}_{a'})}{f(\hat{h}_a)f(\hat{h}_{a'})} \right\}.$$
 (1)

# Available Key bits $(I_{\rm K})$

- Mutual Information can also be computed from Entropy.
- Entropy: Average information gained by observing a single variable

$$H(\hat{h}_a) = \int f(\hat{h}_a) \log f(\hat{h}_a) d\hat{h}_a.$$
 (2)

 Joint Entropy: Average total information gained by observing two or more variables

$$H(\hat{h}_{a}, \hat{h}_{a'}) = \int f(\hat{h}_{a}, \hat{h}_{a'}) \log f(\hat{h}_{a}, \hat{h}_{a'}) d\hat{h}_{a} d\hat{h}_{a'}.$$
 (3)

# Available Key bits $(I_{\rm K})$

Mutual Information in terms of entropy is

$$I_{\mathrm{K}} = \mathrm{H}(\hat{h}_a) + \mathrm{H}(\hat{h}_{a'}) - \mathrm{H}(\hat{h}_a, \hat{h}_{a'}).$$



## Wireless Channel

#### Estimated channels at Alice and Bob are

$$\hat{h}_a = \mathbf{E}_{\mathbf{A}}(\theta', \phi') \alpha \mathbf{E}_{\mathbf{B}}(\theta, \phi) + \epsilon_a.$$
(4)

$$\hat{h}_{a'} = \mathbf{E}_{\mathrm{B}}(\theta, \phi) \alpha \mathbf{E}_{\mathrm{A}}(\theta', \phi') + \epsilon_{a'}.$$
(5)



# Wireless Channel using Reconfigurable Antenna

- Alice has a reconfigurable antenna.
- Each reconfigurable element (RE) is a switch.



- Channel is generated by changing the states of REs using an i.i.d uniform distribution.
- Channel distribution is unknown.

Wireless Channel using Reconfigurable Antenna

•  $x_1 = h_a$  for  $N_{RE} = 24$ ,  $x_2 = h_a$  for  $N_{RE} = 8$ .



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# Available Key bits Computation

#### Approaches

- Gaussian Approximation  $(I_{
  m K,GA})$
- Numerical Computation  $(I_{
  m K,NC})$
- Histogram based Approximation  $(I_{
  m K,HA})$
- Gaussian Mixtures based Approximation  $(I_{
  m K,GM})$
- Channels
  - By assuming  $h_a$  is Gaussian
  - By computing  $h_a$  for  $N_{\rm RE} = 24$
  - By computing  $h_a$  for  $N_{\rm RE} = 8$

## Gaussian Approximation

Closed form solution of entropy exists [1]

$$I_{\rm K,GA} = {\rm H}(\hat{h}_a) + {\rm H}(\hat{h}_{a'}) - {\rm H}(\hat{h}_a, \hat{h}_{a'}),$$

$$= \log_2(\pi e)\sigma_{\hat{h}_a}^2 + \log_2(\pi e)\sigma_{\hat{h}_{a'}}^2 - \log_2(\pi e)^2 |\hat{\mathbf{R}}_{h_a h_{a'}}|.$$
(6)

Channel	$I_{ m K,GA}$ (in bits)	
Gaussian	2.5266	
$N_{\rm RE} = 24$	2.3074	
$N_{\rm RE} = 8$	2.0643	

## Numerical Computation

Mutual information needs to be computed numerically [2]

$$I_{\rm K,NC} = I(\hat{h}_a; \hat{h}_{a'}) = \mathbf{E} \left\{ \log_2 \frac{f(\hat{h}_a, \hat{h}_{a'})}{f(\hat{h}_a)f(\hat{h}_{a'})} \right\}.$$
 (7)

► The individual pdfs f(ĥ<sub>a</sub>) and f(ĥ<sub>a'</sub>) can be expressed in terms of the conditional pdfs

$$f(\hat{h}_a) = \int f(\hat{h}_a|h_a)dh_a = \mathcal{E}_{h_a}f(\hat{h}_a|h_a), \qquad (8)$$
$$= \mathcal{E}_{h_a}f_n((\hat{h}_a - h_a)/\sigma_a^2).$$

Similarly,

$$f(\hat{h}_{a'}) = \mathcal{E}_{h_a} f_n((\hat{h}_{a'} - h_a) / \sigma_{a'}^2).$$
(9)

#### Numerical Computation

▶ Joint pdf  $f(\hat{h}_a, \hat{h}_{a'})$  can be expressed as the product of two noise pdfs

$$f(\hat{h}_{a}, \hat{h}_{a'}) = \int f(\hat{h}_{a}, \hat{h}_{a'}|h_{a}) dh_{a}$$
(10)  
=  $E_{h_{a}}f(\hat{h}_{a}, \hat{h}_{a'}|h_{a}) = E_{h_{a}}\{f(\hat{h}_{a}|h_{a})f(\hat{h}_{a'}|h_{a})\}$   
=  $E_{h_{a}}\{f_{n}((\hat{h}_{a} - h_{a})/\sigma_{a}^{2})f_{n}((\hat{h}_{a'} - h_{a})/\sigma_{a'}^{2}\}.$ 

The convergence of the numerical computation will depend on N and M, which are the number of sample points in the outer and inner expectations.

### Numerical Computation



### Histogram based Approximation

 Mutual information is computed by estimating the pdfs using multi-dimensional histograms

$$I_{\rm K,HA} = I(\hat{h}_a; \hat{h}_{a'}) = \mathbf{E} \left\{ \log_2 \frac{f(\hat{h}_a, \hat{h}_{a'})}{f(\hat{h}_a)f(\hat{h}_{a'})} \right\}.$$
 (11)

 Estimated channel pdfs can be expressed in terms of convolution as

$$f(\hat{h}_a) = f(h_a) * f(\epsilon_a).$$
(12)

$$f(\hat{h}_{a'}) = f(h_{a'}) * f(\epsilon_{a'}).$$
(13)

## Histogram based Approximation

• The joint pdf  $f(\hat{h}_a, \hat{h}_{a'})$  can be expressed as

$$f(\hat{h}_{a}, \hat{h}_{a'}) = f(h_{a}, h_{a'}) * f(\epsilon_{a}, \epsilon_{a'}).$$
(14)

- 2-D and 4-D convolution and histograms are used for computing individual and joint pdfs.
- Number of bins  $(N_B)$  along each dimension is variable.

## Histogram based Approximation



#### Gaussian Mixture based Approximation

 A given distribution can be expressed in terms of a mixture of several Gaussian distributions

$$\hat{f}(\angle x) = \sum_{i=1}^{L} w_i \cdot \mathcal{N}(\angle x, \mu_i, \sigma_i^2).$$
(15)



### Gaussian Entropy Computation

Mutual information is computed using the entropy as

$$I_{\rm K,GM} = {\rm H}(\hat{h}_a) + {\rm H}(\hat{h}_{a'}) - {\rm H}(\hat{h}_a, \hat{h}_{a'}).$$
(16)

► The entropy for a random vector <u>x</u> of size P with pdf f(<u>x</u>) is given by

$$H(\underline{x}) = \int f(\underline{x}) \cdot \log f(\underline{x}) d\underline{x}.$$
 (17)

where

$$f(\underline{x}) = \sum_{i=1}^{L} w_i \cdot \mathcal{N}(\underline{x}, \underline{\mu}_i, \mathbf{C}_i).$$
(18)

•  $\log f(\underline{x})$  is estimated using Taylor series [3]

$$\log f(\underline{x}) = \sum_{k=0}^{R} \frac{1}{k!} \left( \left( \underline{x} - \underline{\mu_i} \right) \odot \Delta \right)^k \log f(\underline{x}) |_{\underline{x} = \underline{\mu_i}} + O_R.$$
 (19)

### Gaussian Entropy Computation

For Gaussian distribution  $I_{K,GM} = 2.5216$  when L = 1.



# Conclusion

- Comparison of different techniques for arbitrary channel distribution.
- ▶ Case A:  $h_a$  is generated using  $N_{\rm RE} = 24$
- Case B:  $h_a$  is generated using  $N_{\rm RE} = 8$

Method Used	$I_{ m K}$ for Case A	$I_{ m K}$ for Case B	Time
Gaussian Approx	2.3074	2.0643	$\leq 1$
Numerical Computation	2.2737	1.9142	1920
Histogram based Approx	2.2581	1.8908	13289
Gaussian mixtures	2.2787	1.9106	209

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